

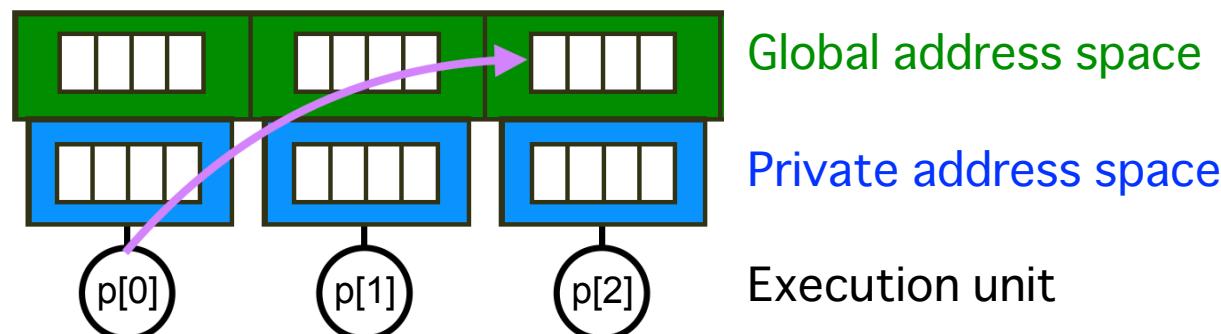
Performance Evaluation for Omni XcalableMP Compiler on Many-core Cluster System based on Knights Landing

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- 2.Center for Computational Sciences University of Tsukuba

Background (1/2)

- Partitioned Global Address Space (PGAS) programming model for cluster system
 - Provide **global address space** on distributed memory system
 - Higher productivity than MPI
 - **XcalableMP (XMP)**, XcalableACC, DASH, Coarray Fortran, Unified Parallel C (UPC), UPC++, X10, Chapel and so on



Background (2/2)

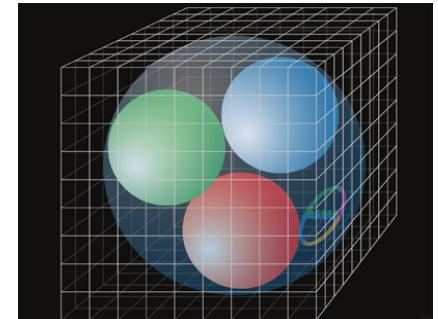
- **XMP** is a directive-based language extension
 - Based on C and Fortran (C++ on the table)
 - Collaborate with OpenMP directives for thread programming
 - Designed by PC cluster consortium
 - <http://xcalablemp.org>
- **Omni compiler**
 - Reference implementation for XMP
 - Developed by RIKEN AICS and University of Tsukuba
 - Source-to-Source compiler
 - Support : The K computer, **Intel Xeon Phi Cluster**, Cray machines, ...
 - <http://omni-compiler.org>

Objective

- Little experience with Omni compiler on Intel Xeon Phi cluster system, which is attracting attention in the HPC field



- Evaluate the performance of Omni compiler on Oakforest-PACS, which is a cluster system based on Knights Landing (9th in the latest Top500 list)
- Make the following key contributions:
 - Evaluation of the Lattice QCD mini-application using a hybrid model of XMP and OpenMP on Oakforest-PACS
 - Effective code translation method for a source-to-source compiler



Agenda from this slide

- Overview of XMP and Omni compiler
- Performance tuning of Omni compiler on a single compute node
- Evaluation of the Lattice QCD mini-application on Oakforest-PACS
- Summary

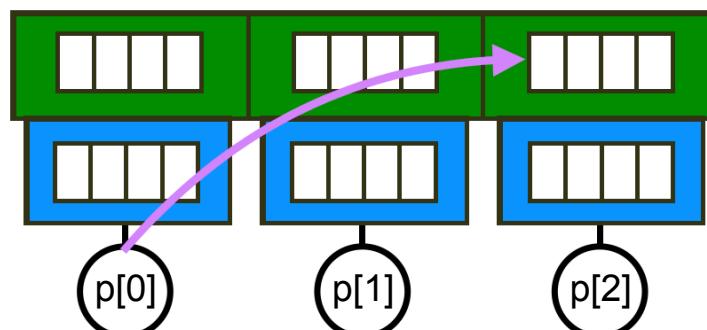
Example of XcalableMP programming

```
int a[MAX];  
#pragma xmp nodes p[3]  
#pragma xmp template t[MAX]  
#pragma xmp distribute t[block] onto p  
#pragma xmp align a[i] with t[i]  
  
int main(){  
#pragma xmp loop on t[i]  
    for(int i = 0; i < MAX; i++)  
        a[i] = foo(i);
```

Define execution unit and data distribution

Parallelize loop statement

1. XMP loop directive parallelizes across execution units



Global address space

Private address space

Execution unit

Example of XcalableMP programming

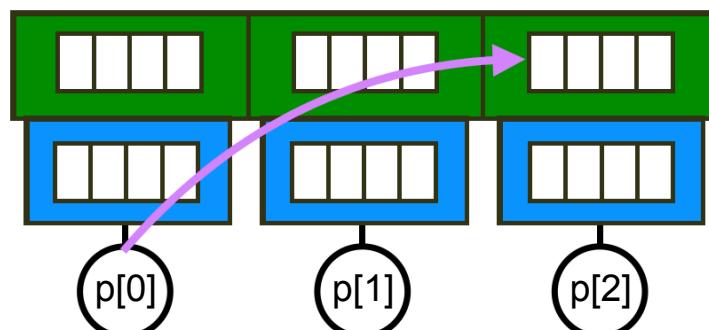
```
int a[MAX];  
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```

Define execution unit and data distribution

```
int main(){  
#pragma xmp loop on t[i]  
#pragma omp parallel for  
    for(int i = 0; i < MAX; i++)  
        a[i] = foo(i);
```

Parallelize loop statement

1. **XMP loop** directive parallelizes across execution units
2. **OpenMP parallel for** directive parallelizes across threads



Global address space

Private address space

Execution unit

Example of XcalableMP programming

```
int a[MAX];
```

```
#pragma xmp nodes p[*]
```

```
#pragma xmp template t[MAX]
```

```
#pragma xmp distribute t[block] onto p
```

```
#pragma xmp align a[i] with t[i]
```

Number of execution units is defined dynamically

Define execution unit and
data distribution

```
int main(){
```

```
#pragma xmp loop on t[i]
```

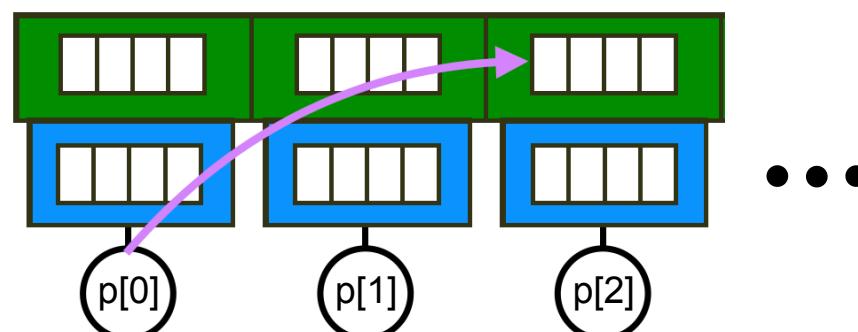
```
#pragma omp parallel for
```

```
for(int i = 0; i < MAX; i++)
```

```
    a[i] = foo(i);
```

Parallelize loop statement

1. XMP loop directive parallelizes across execution units
2. OpenMP parallel for directive parallelizes across threads

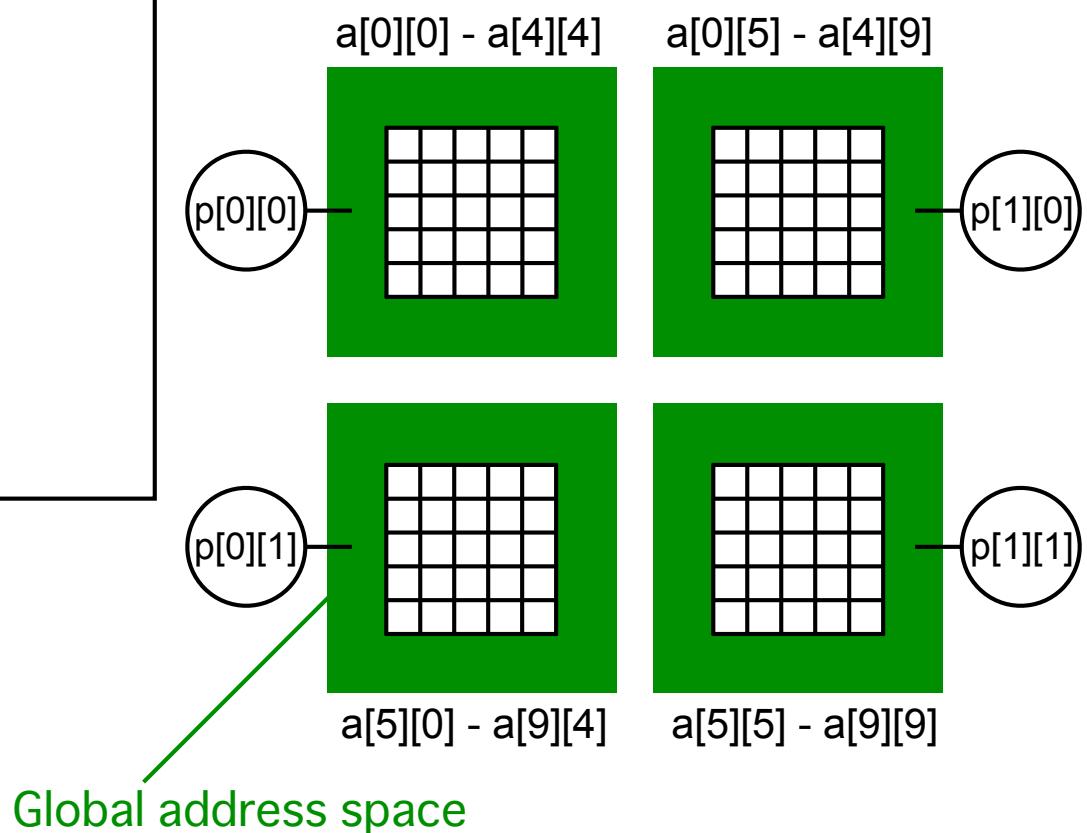


Example of XcalableMP programming

- Declaration of multi-dimensional array on global memory address

```
int a[10][10];
#pragma xmp nodes p[2][2]
#pragma xmp template t[10][10]
#pragma xmp distribute t[block][block] onto p
#pragma xmp align a[i][j] with t[i][j]

int main(){
#pragma xmp loop (i,j) on t[i][j]
#pragma omp parallel for collapse(2)
    for(int i = 0; i < 10; i++)
        for(int j = 0; j < 10; j++)
            a[i][j] = foo(i,j);
```



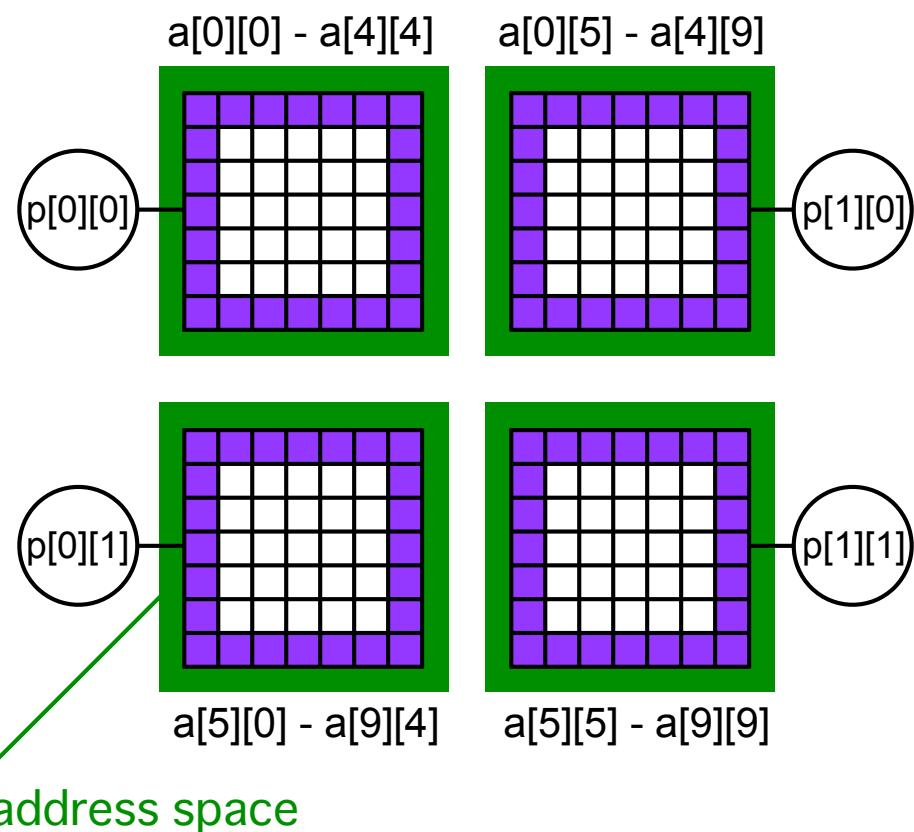
Example of XcalableMP programming

- Shadow/Reflect directives for Stencil application

```
int a[10][10];
#pragma xmp nodes p[2][2]
#pragma xmp template t[10][10]
#pragma xmp distribute t[block][block] onto p
#pragma xmp align a[i][j] with t[i][j]
#pragma xmp shadow a[1][1]

int main(){
    :
    #pragma xmp reflect (a) width(/periodic/1) ¥
        orthogonal
    #pragma xmp loop (i,j) on t[i][j]
    #pragma omp parallel for
    for(int i = 0; i < 10; i++)
        for(int j = 0; j < 10; j++)
            ... = a[i-1][j] + a[i+1][j] + a[i][j-1] + a[i][j+1];
```

Shadow directive is to add halo region in distributed array



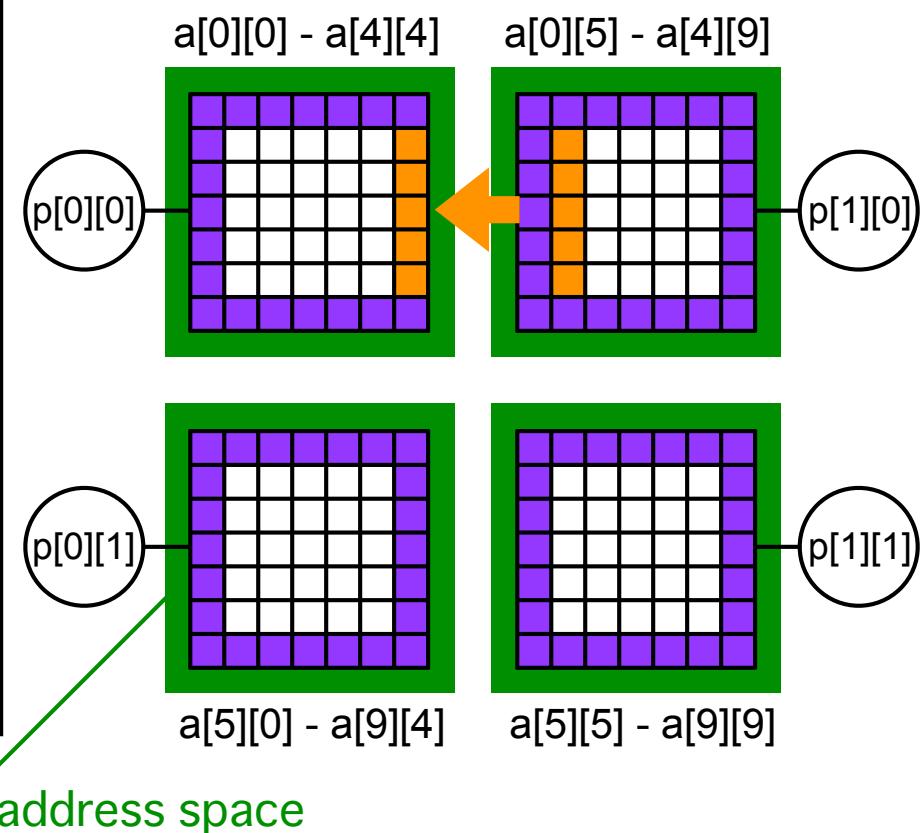
Example of XcalableMP programming

- Shadow/Reflect directives for Stencil application

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int main(){
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    #pragma xmp loop (i,j) on t[i][j]
    #pragma omp parallel for
    for(int i = 0; i < 10; i++)
        for(int j = 0; j < 10; j++)
            ... = a[i-1][j] + a[i+1][j] + a[i][j-1] + a[i][j+1];
```

Reflect directive is to exchange halo region among neighborhood nodes



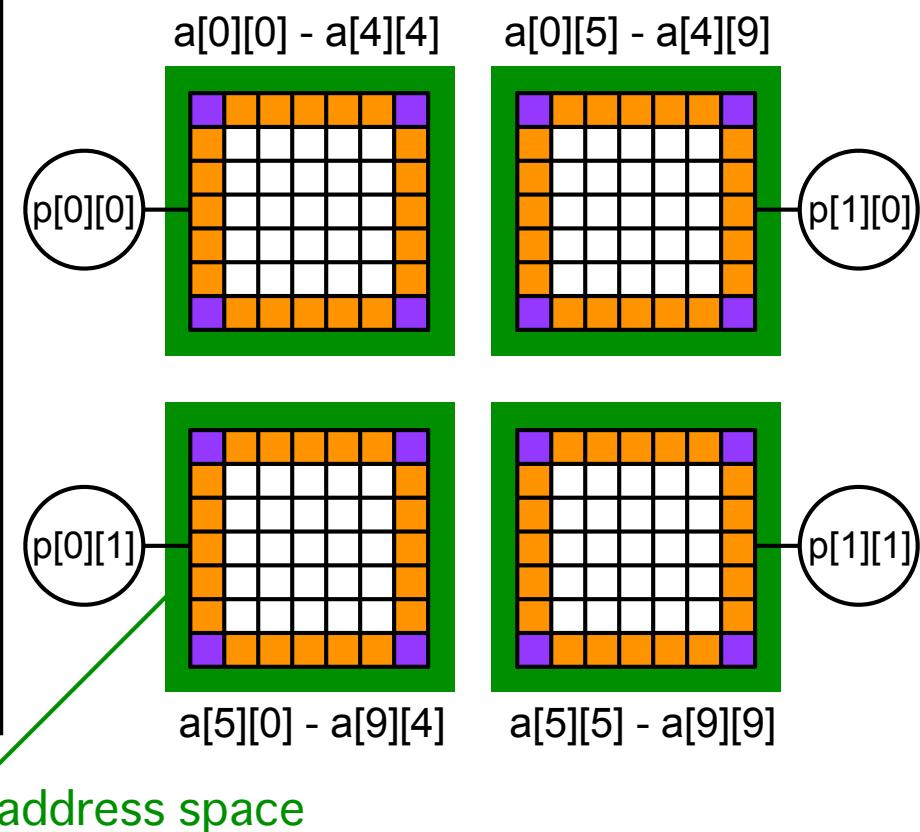
Example of XcalableMP programming

- Shadow/Reflect directives for Stencil application

```
int a[10][10];
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        for(int j = 0; j < 10; j++)
            ... = a[i-1][j] + a[i+1][j] + a[i][j-1] + a[i][j+1];
```

Reflect directive is to exchange halo region among neighborhood nodes



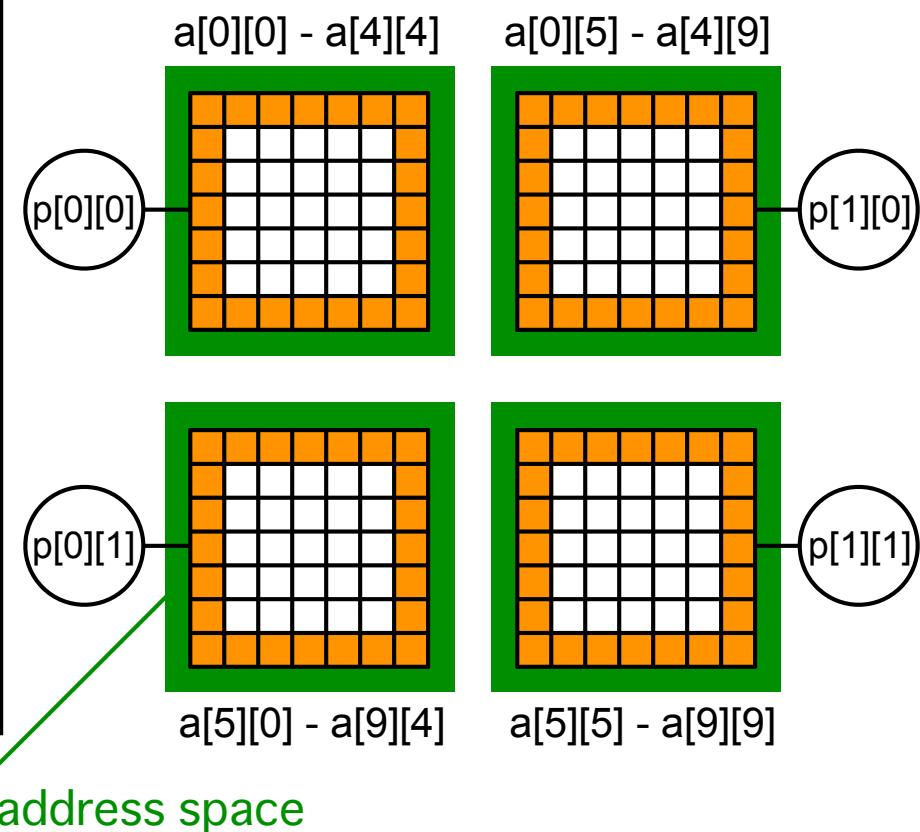
Example of XcalableMP programming

- Shadow/Reflect directives for Stencil application

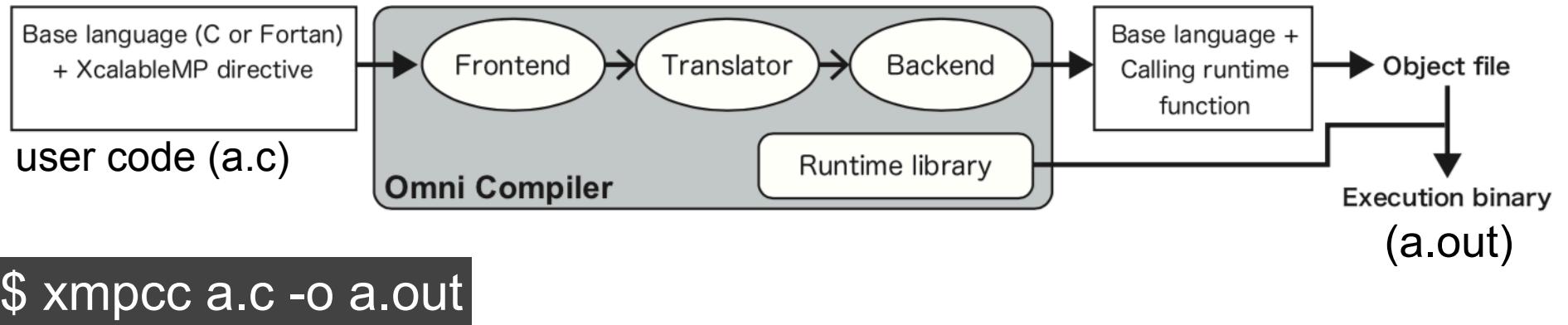
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    for(int i = 0; i < 10; i++)
        for(int j = 0; j < 10; j++)
            ... = a[i-1][j] + a[i+1][j] + a[i][j-1] + a[i][j+1];
```

Reflect directive is to exchange halo region among neighborhood nodes



Omni compiler



- **Source-to-Source compiler**

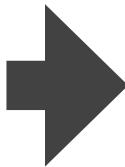
- A user code with XMP directives is translated to a parallel code with runtime calls of Omni compiler's runtime library
- The Omni compiler's runtime library is implemented in C and MPI
- The translated parallel code is compiled by a native compiler
 - e.g. GNU, Intel, PGI, Cray, and so on

Examples of code translation

- Define distributed array on global address space

User code

```
double a[10][10];  
#pragma xmp align a[i][j] with t[i][j]
```



Translated code

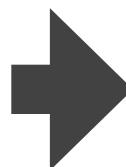
```
int *_XMP_ADDR_a;  
:  
_XMP_alloc_array(&_XMP_ADDR_a, ...);
```

Multi-dimensional array is expressed as a pointer and memory is allocated dynamically. One of reasons why memory is allocated dynamically, memory size may be defined dynamically (e.g. using `#pragma xmp nodes p[*]`).

- Parallelize loop statement on 2x2 execution units

User code

```
#pragma xmp loop (i,j) on t[i][j]  
for(int i=0;i<10;i++)  
    for(int j=0;j<10;j++)  
        a[i][j] = ...;
```



Translated code

```
for(int i=0;i<5;i++)  
    for(int j=0;j<5;j++)  
        *(_XMP_ADDR_a + i*5 + j) = ..;
```

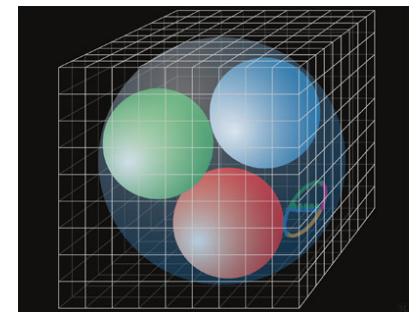
- The initial value and ending condition use constants automatically as possible
- An array operation is translated to a pointer operation

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- Overview of XMP and Omni compiler
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Overview of Lattice QCD

- Lattice QCD is a discrete formulation of QCD (Quantum Chromodynamics)
 - Describe the strong interaction among “quarks” and “gluons”
 - Quark is a species of elementary particles
 - Gluon is a particle that works between quarks
- Lattice QCD is formulated on a four-dimensional lattice (Time and XYZ axes)
- Our Lattice QCD code is based on an existing Lattice QCD mini-application (<http://research.kek.jp/people/matufuru/Research/Programs/>)
 - By High Energy Accelerator Research Organization, Japan
 - Implemented by extracting the main kernel of the Bridge++, which is a real-world application for lattice gauge theories including QCD (http://bridge.kek.jp/Lattice-code/index_e.html)



Overview of algorithm

- Pseudo-code (CG method is used)

```
S = B          // COPY
R = B          // COPY
X = B          // COPY
sr = norm(S)  // NORM
T = WD(U,X)   // Main Kernel
S = WD(U,T)   // Main Kernel
R = R - S     // AXPY
P = R          // COPY
rrp = rr = norm(R) // NORM
do{
    T = WD(U,P)   // Main Kernel
    V = WD(U,T)   // Main Kernel
    pap = dot(V,P) // DOT
    cr = rr/pap
    X = cr * P + X // AXPY
    R = -cr * V + R // AXPY
    rr = norm(R)   // NORM
    bk = rr/rrp
    P = bk * P     // SCAL
    P = P + R      // AXPY
    rrp = rr
}while(rr/sr > 1.E-16)
```

WD() is the Wilson-Dirac operator

$$D_{x,y} = \delta_{x,y} - \kappa \sum_{\mu=1}^4 \{(1 - \gamma_\mu)U_\mu(x)\delta_{x+\hat{\mu},y} + (1 + \gamma_\mu)U_\mu^\dagger(x - \hat{\mu})\delta_{x-\hat{\mu},y}\}$$

- Main kernel (most costly)
- Stencil calculation

#pragma xmp reflect (X) width(..) orthogonal

WD(X, ...);

void WD(Quark_t X[NT][NZ][NY][NX], ...){

:

#pragma xmp loop (t,z) on t[t][z]

#pragma omp parallel for collapse(4)

for(int t=0;t<NT;t++)

for(int z=0;z<NZ;z++)

for(int y=0;y<NY;y++)

for(int x=0;x<NX;x++){

:

Condition of Preliminary Evaluation

- Oakforest-PACS as a KNL cluster system

CPU	Intel Xeon Phi 7250 1.4–1.6GHz 68Cores
Memory	MCDRAM 16GB, DDR4 96GB
Network	Intel Omni-Path Host Fabric Interface 12.5GB/s
Software	intel/2017.4.196, intelmpi/2017.3.196



- COMA as a general PC cluster system

CPU	Intel Xeon-E5 2670v2 2.5–3.3GHz 10Cores, 2Sockets
Memory	DDR3 64GB
Network	InfiniBand FDR 7GB/s
Software	intel/17.0.5, intelmpi/2017.4

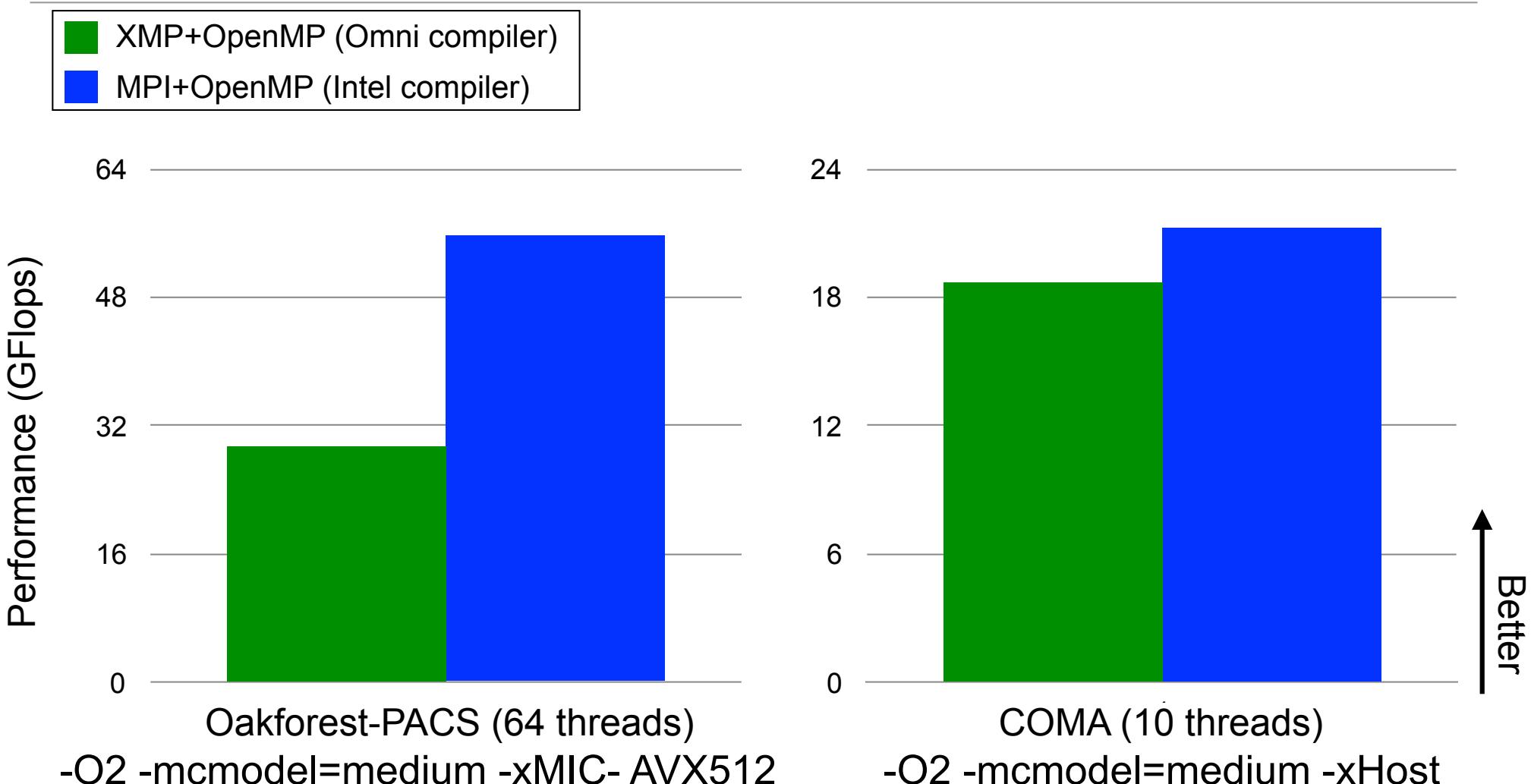
Note: Although COMA has KNC, we didn't use it.

For comparison purpose, we also developed Lattice QCD mini-application in MPI+OpenMP.

one process with multi-threads
in a single compute node



Result of Preliminary Evaluation (NT, NZ, NY, NX) = (32,32,32,32)



Although Omni Compiler also uses Intel compiler as a backend compiler, performance results of **XMP+OpenMP** are worse than those of **MPI+OpenMP**.

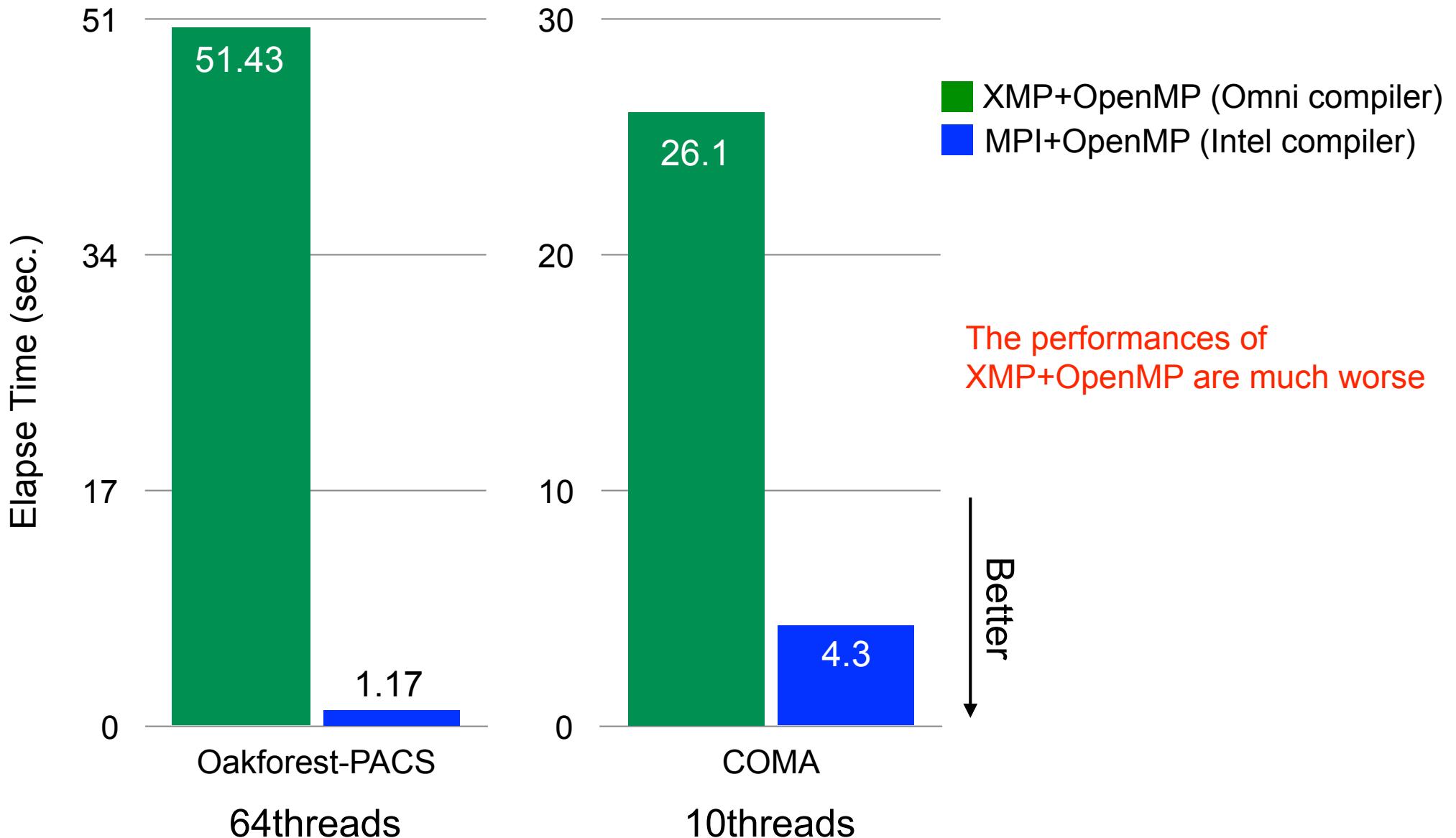
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S = WD(U,T)   // Main Kernel
R = R - S    // AXPY
P = R          // COPY
rrp = rr = norm(R) // NORM
do{
    T = WD(U,P)   // Main Kernel
    V = WD(U,T)   // Main Kernel
    pap = dot(V,P) // DOT
    cr = rr/pap
    X = cr * P + X // AXPY
    R = -cr * V + R // AXPY
    rr = norm(R)   // NORM
    bk = rr/rrp
    P = bk * P    // SCAL
    P = P + R    // AXPY
    rrp = rr
}while(rr/sr > 1.E-16)
```

We profiled all functions.
As a result, we found that the performance results of **mathematical functions** were worse.
Especially, the performance result of **SCAL** was much worse.

Result of SCAL (NT, NZ, NY, NX) = (32,32,32,32)



SCAL function (A part of Lattice QCD code)

User code

```
typedef struct Quark {  
    double v[4][3][2];  
} Quark_t;  
Quark_t X[NT][NZ][NY][NX];  
#pragma xmp nodes p[1][1]  
:  
#pragma xmp shadow (X[1][1][0][0])
```

Multiply the given vector by the given scalar
($X := a * X$)

```
void scal(Quark_t X[NT][NZ][NY][NX],  
          const double a){  
    :  
#pragma xmp loop (t,z) on t[t][z]  
#pragma omp parallel for collapse(4)  
for(int t=0;t<NT;t++)  
    for(int z=0;z<NZ;z++)  
        for(int y=0;y<NY;y++)  
            for(int x=0;x<NX;x++)  
                for(int i=0;i<4;i++)  
                    for(int j=0;j<3;j++)  
                        for(int k=0;k<2;k++)  
                            X[t][z][y][x].v[i][j][k] *= a;
```

Tuning Omni compiler

Translated code (old)

```
void scal(Quark_t *_XMP_ADDR_X,
          const double a){
    :
#pragma omp parallel for collapse(4)
for(int t=0;t<NT;t++)
    for(int z=0;z<NZ;z++)
        for(int y=0;y<NY;y++)
            for(int x=0;x<NX;x++)
                for(int i=0;i<4;i++)
                    for(int j=0;j<3;j++)
                        for(int k=0;k<2;k++)
                            (*(((*(((_XMP_ADDR_X +
(t+1)*(NZ+2)*(NY)*(NX) +
(z+1)*(NY)*(NX) + y*(NX) + x
->v) + i)) + j)) + k)) *= a;
```

Translated code (new)

```
void scal(Quark_t *_XMP_ADDR_X,
          const double a){
    :
Quark_t (*X_NEW)[NZ+2][NY][NX] =
    (Quark_t (*)[NZ+2][NY][NX])_XMP_ADDR_X;
# pragma omp parallel for collapse(4)
for(int t=0;t<NT;t++)
    for(int z=0;z<NZ;z++)
        for(int y=0;y<NY;y++)
            for(int x=0;x<NX;x++)
                for(int i=0;i<4;i++)
                    for(int j=0;j<3;j++)
                        for(int k=0;k<2;k++)
                            (*(((*(((*(((&(X_NEW[t+1][z+1][y][x])
->v) + i)) + j)) + k)) *= a;
```

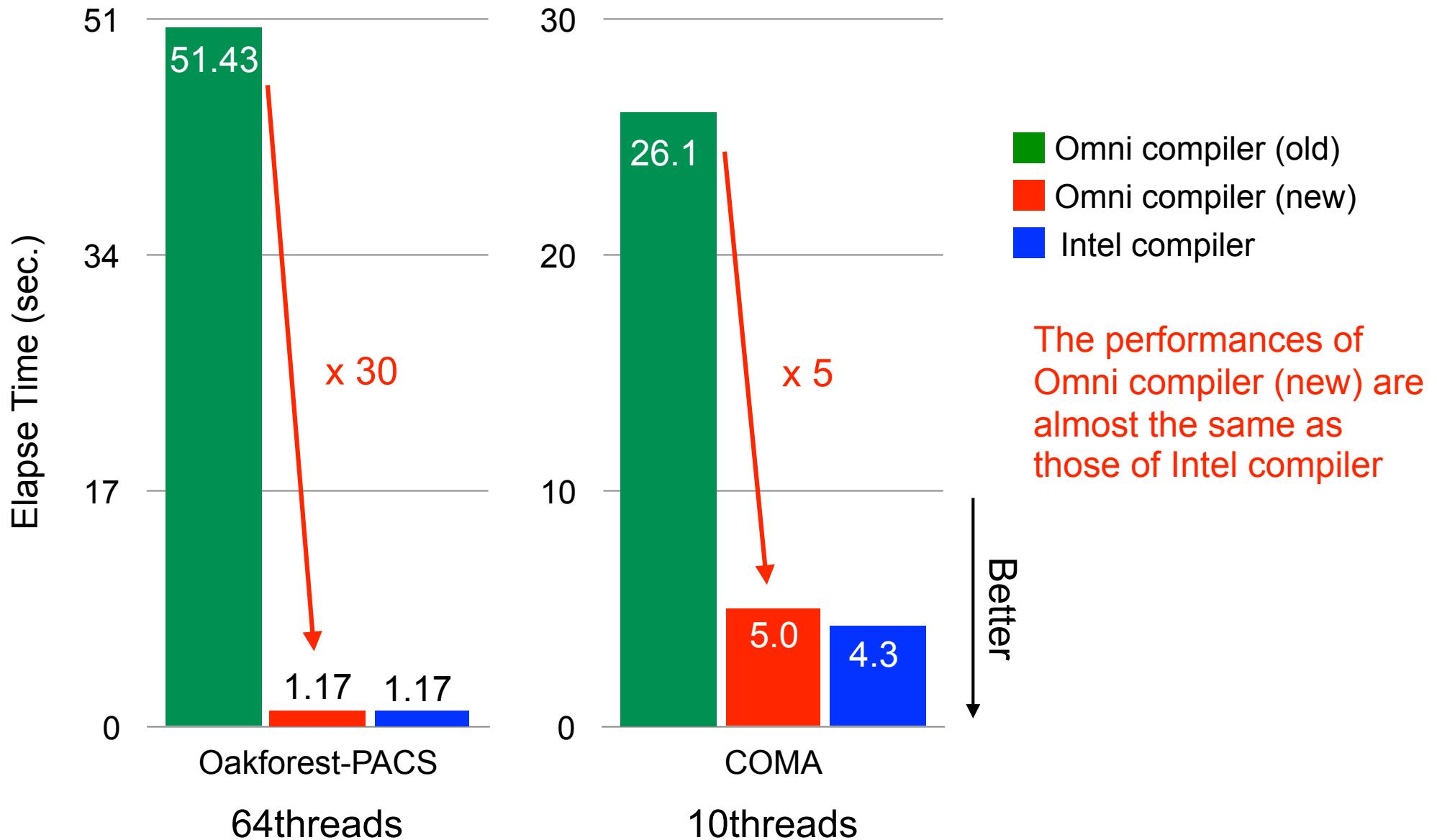
The reason why the performance is worse is that the size of each dimension has disappeared.

We add "-qopt-report" option to Intel compiler.
"LOOP WAS NOT VECTORIZED"

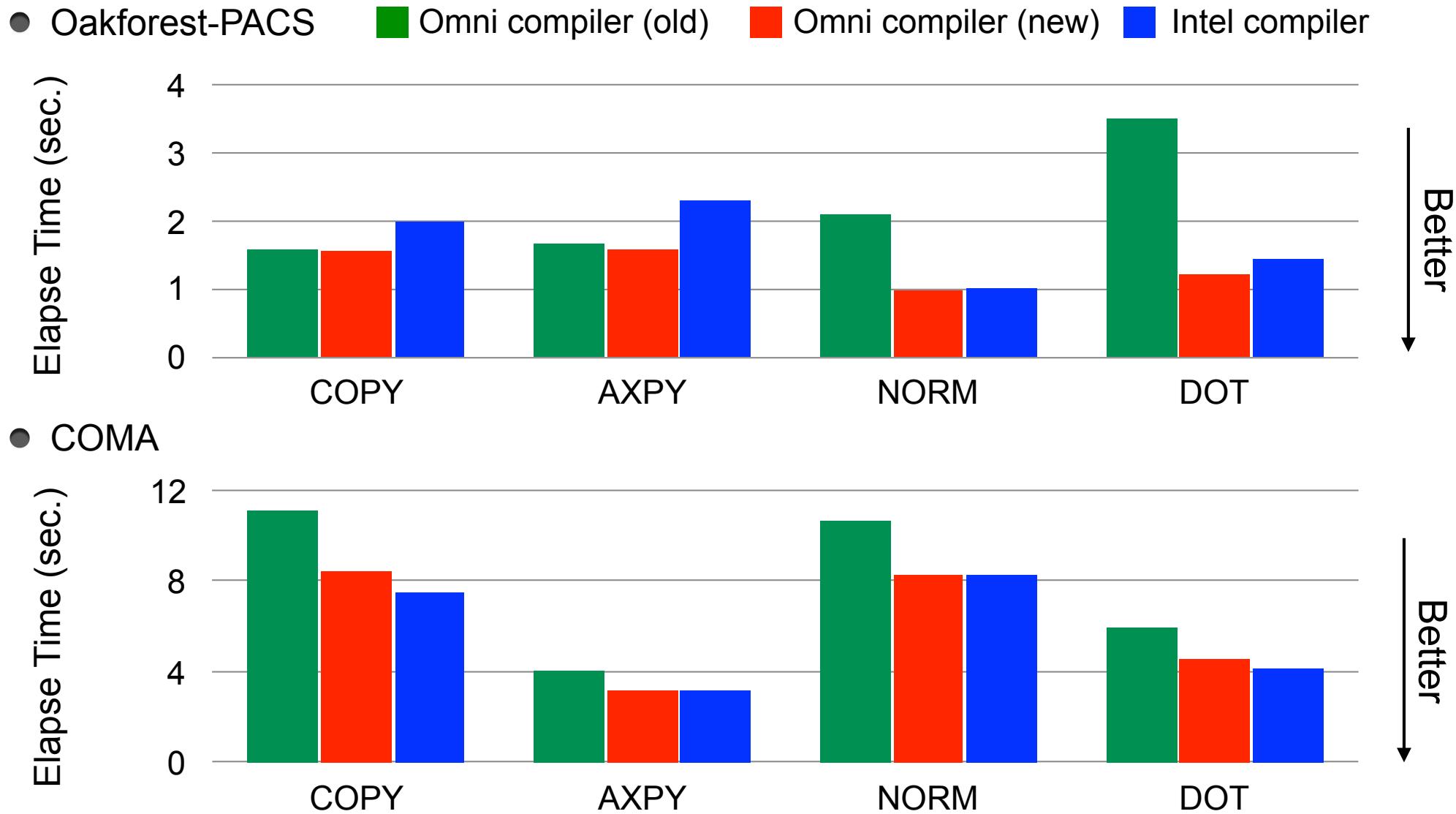
Add a new pointer for a distributed array, which has the size of each dimension.

"LOOP WAS VECTORIZED"

Result of SCAL (NT, NZ, NY, NX) = (32,32,32,32)



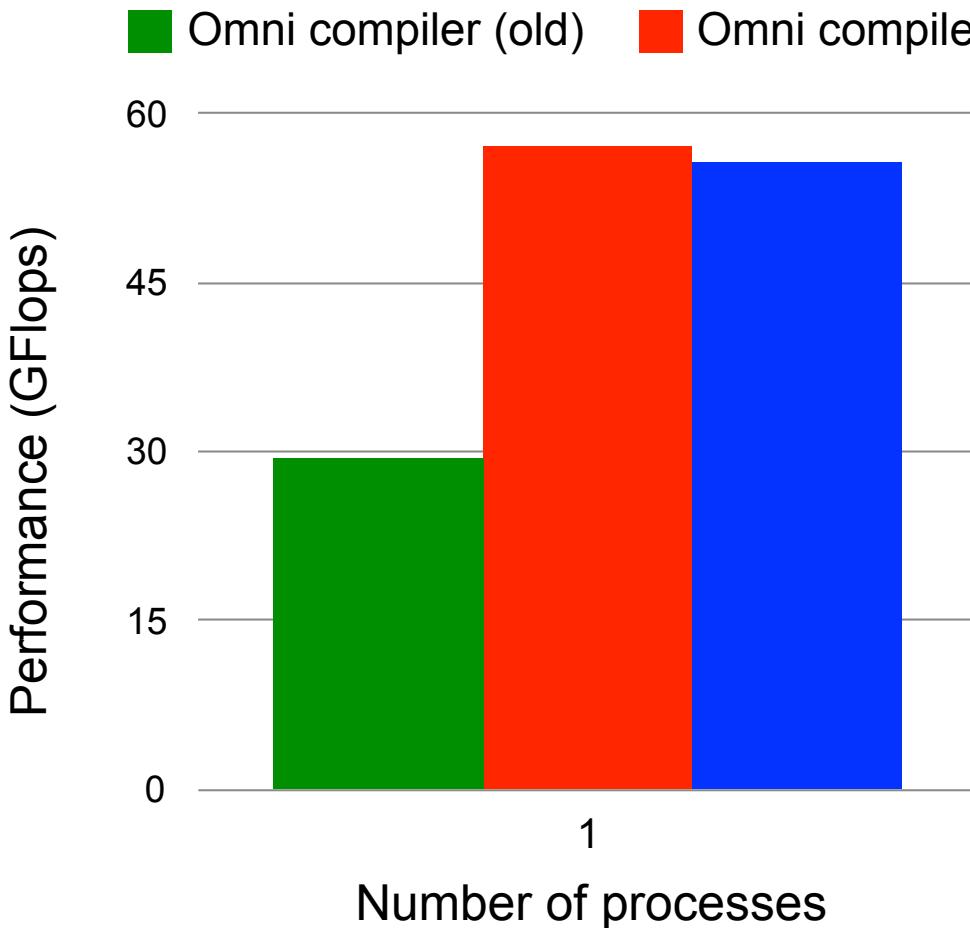
Other mathematical functions (NT, NZ, NY, NX) = (32,32,32,32)



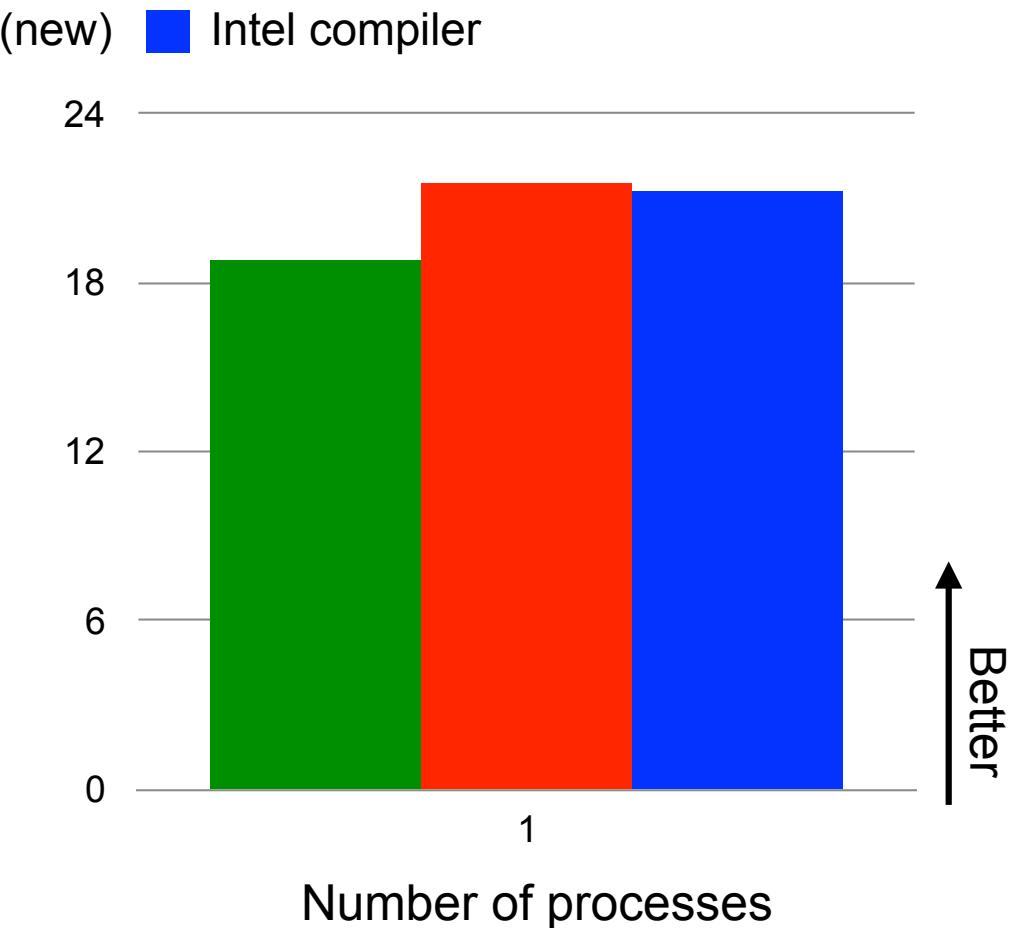
The performances of Omni compiler (new) are better than those of Omni compiler (old)

Result of Preliminary Evaluation (NT, NZ, NY, NX) = (32,32,32,32)

- Oakforest-PACS (64 threads)



- COMA (10 threads)



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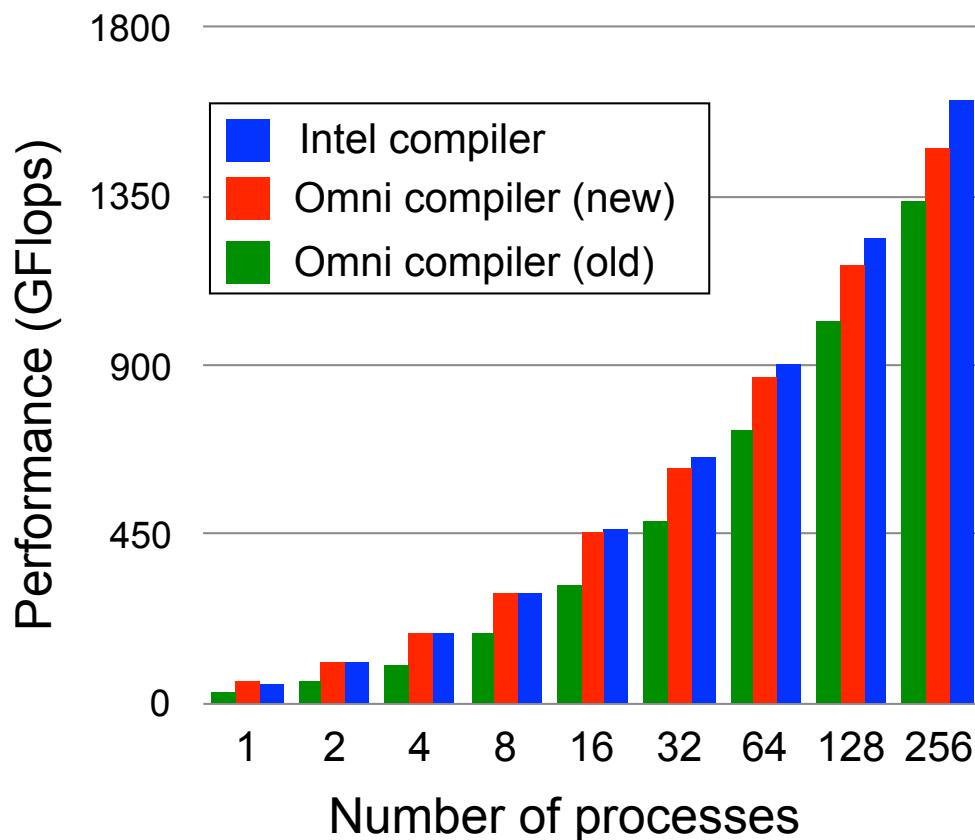
Performance Evaluation on cluster

- Oakforest-PACS and COMA
 - One process per compute node on Oakforest-PACS
 - Two processes per compute node on COMA because it has two CPU sockets
- Problem size is (32,32,32,32) as (NT,NZ,NY,NX) with strong scaling

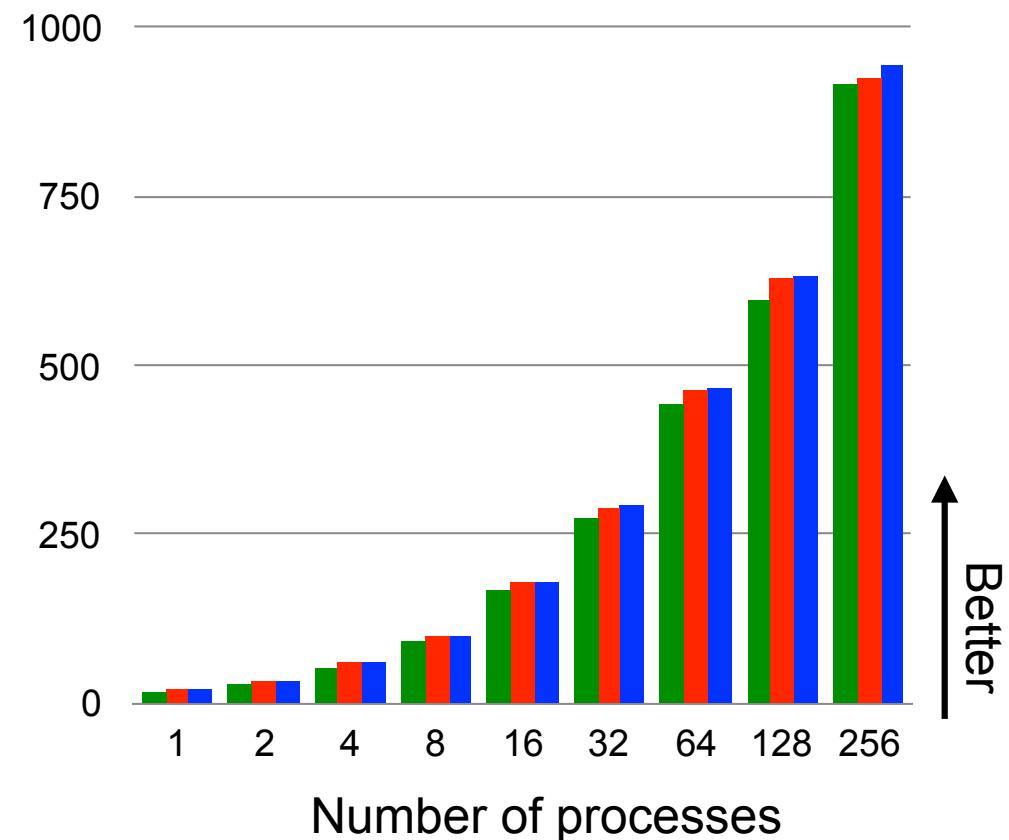


Performance Evaluation

- Oakforest-PACS



- COMA



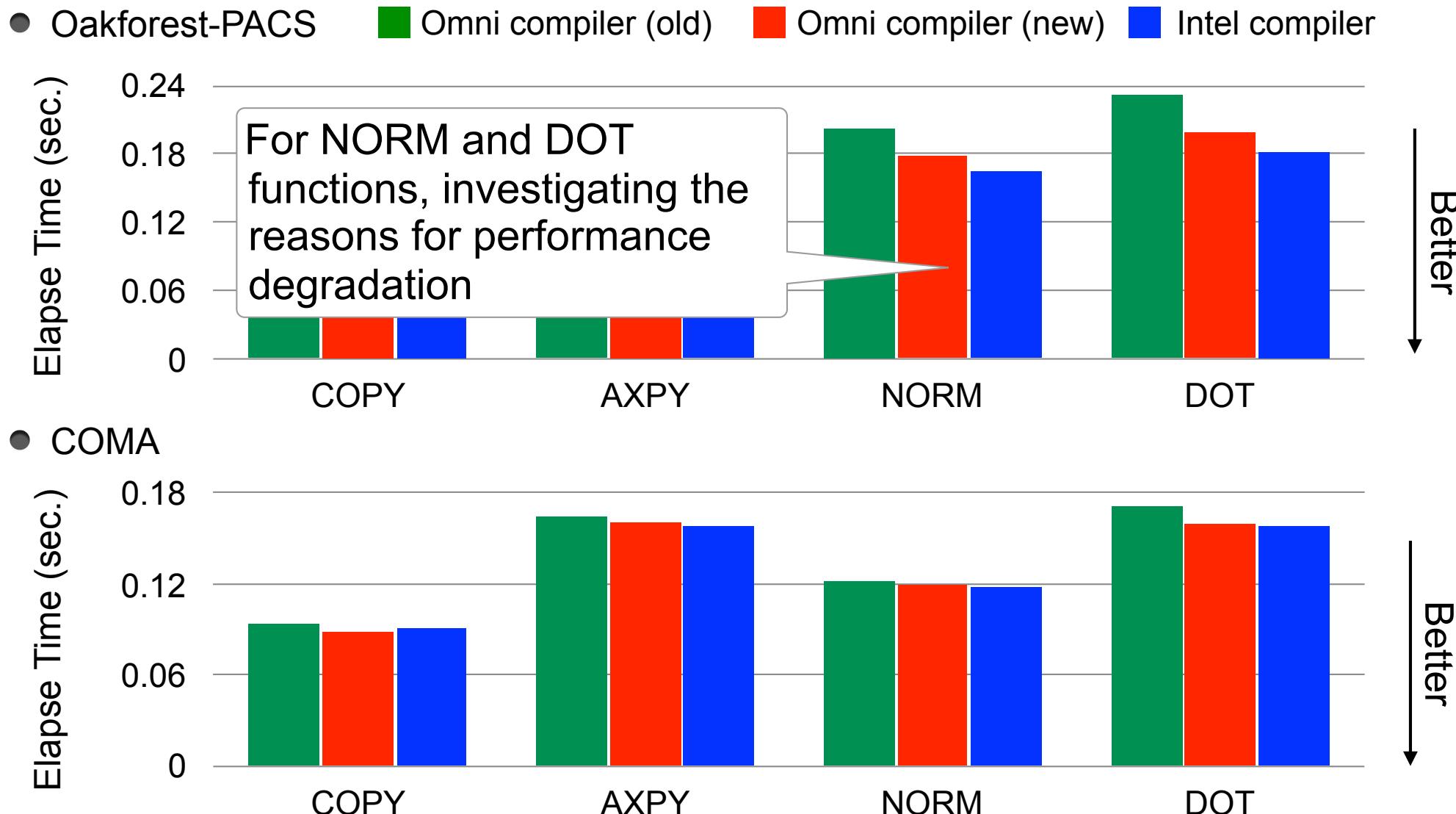
Better ↑

Performances of **Omni compiler (new)** are always better than those of **Omni compiler (old)**.

Performances of **Omni compiler (new)** achieve 94 - 105% of those of **Intel compiler**.

Other mathematical functions (NT, NZ, NY, NX) = (2,2,32,32)

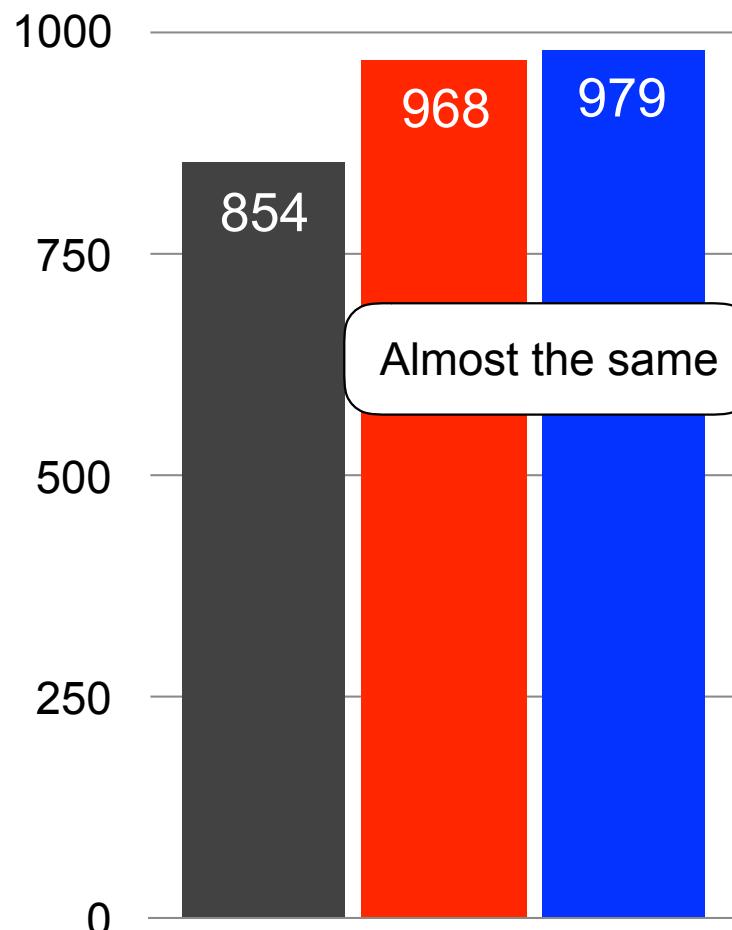
This problem size is a case of using 256 processes.



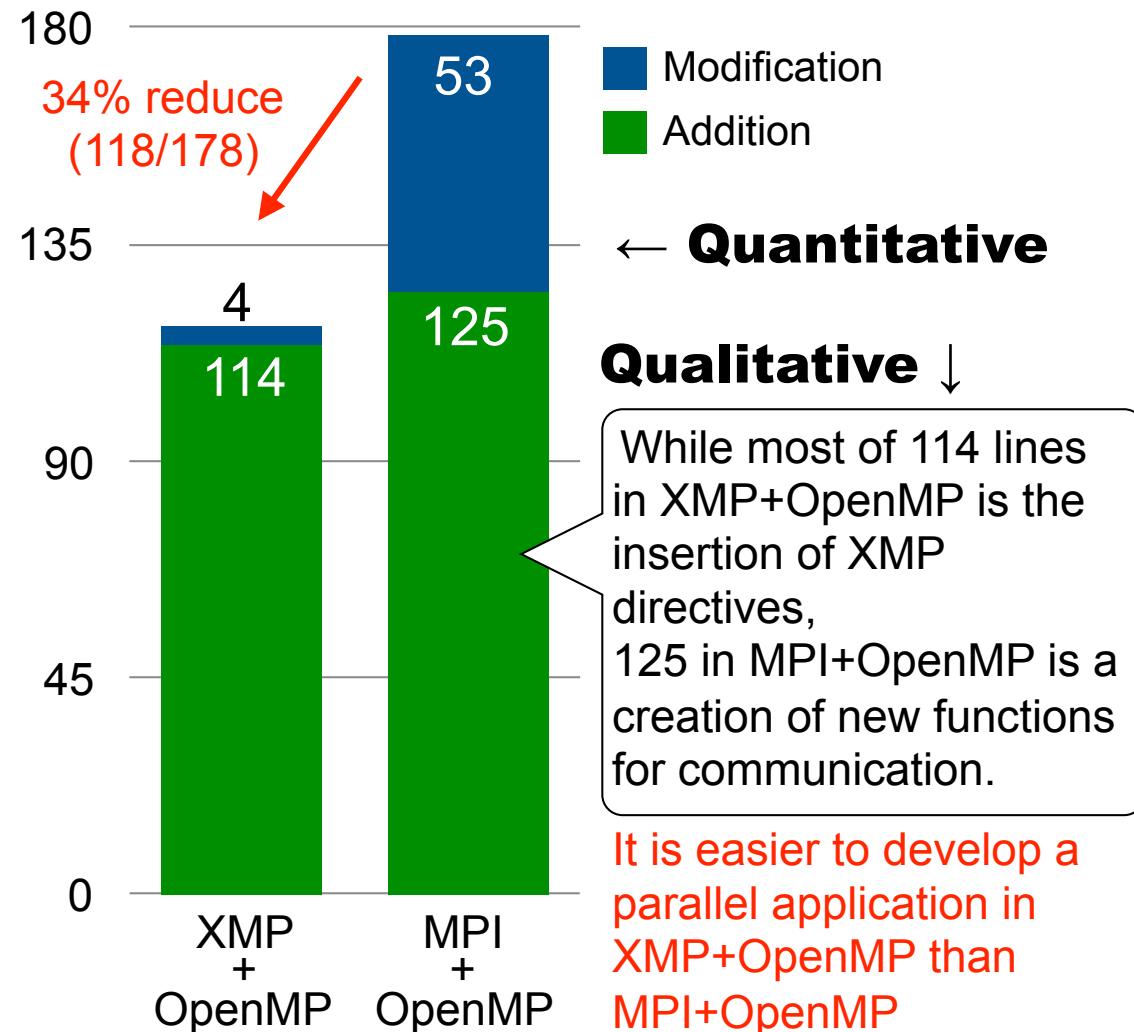
Productivity Evaluation

- Source lines of codes (SLOC)

■ OpenMP (Base code) ■ XMP+OpenMP
■ MPI+OpenMP



- Delta SLOC How many lines the code changed from a base code to a parallel code



Agenda from this slide

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Conclusion

- We evaluated the performance of the Omni compiler on **Oakforest-PACS**, which is a cluster system based on KNL, and **COMA**, which is a general Linux cluster.
- We tuned performance of Omni compiler
 - The SCAL performance of Omni compiler (new) achieves **30 times** better than that of Omni compiler (old) on KNL
- We implemented the Lattice QCD mini-application in XMP+OpenMP
 - The performance in XMP+OpenMP using Omni compiler (new) achieves **94 - 105%** of that in MPI+OpenMP
 - The productivity of XMP+OpenMP is **better** than that of MPI+OpenMP